

# Planning a Journey in an Uncertain Environment

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Modélisation & Simulation*



## Aim of this talk

Flavor of  $\neq$  types of **useful strategies** in stochastic environments.

- ▶ Joint paper<sup>1</sup> with J.-F. Raskin (ULB) and O. Sankur (IRISA, Rennes) [[RRS15b](#)]
- ▶ Full paper available on arXiv: [abs/1411.0835](#)

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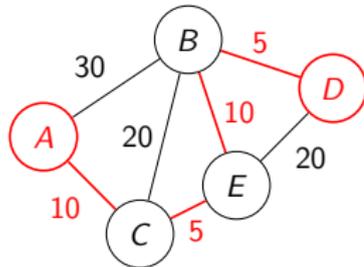
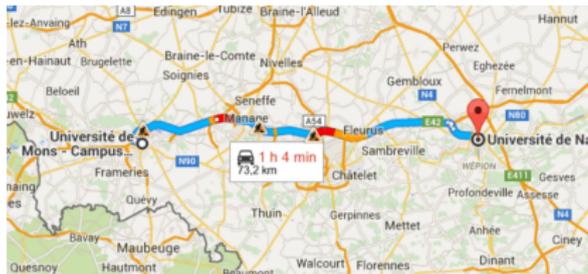
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Applications to the **shortest path problem**.



↪ Find a **path of minimal length** in a weighted graph (Dijkstra, Bellman-Ford, etc) [CGR96].

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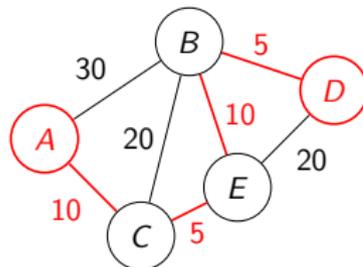
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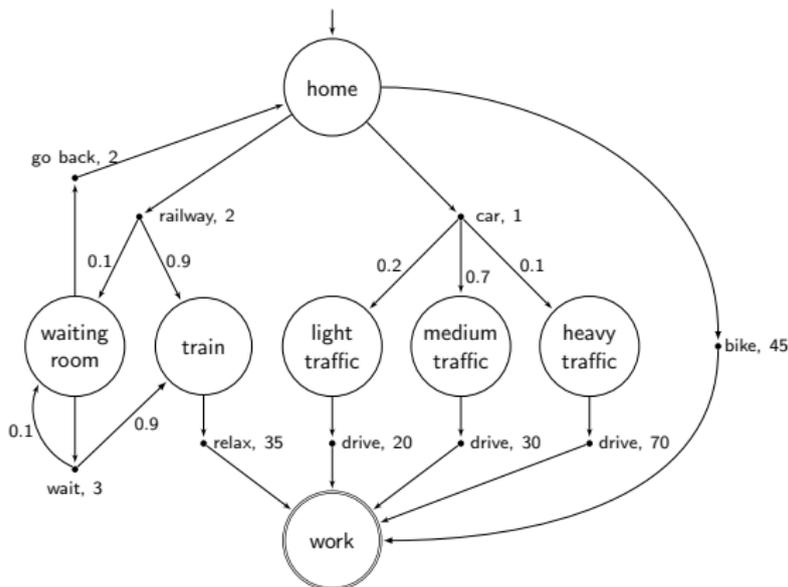


What if the environment is **uncertain**? E.g., in case of heavy traffic, some roads may be crowded.

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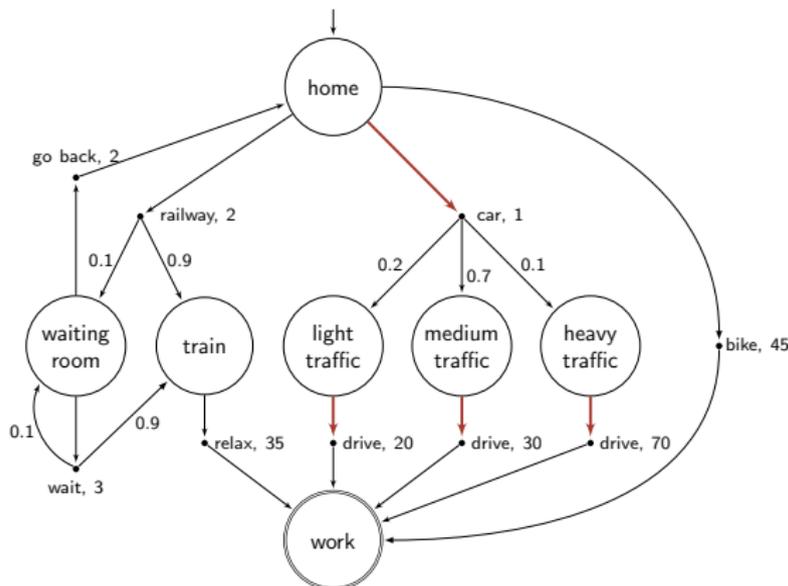
# Planning a journey in an uncertain environment



Each action takes **time**, target = work.

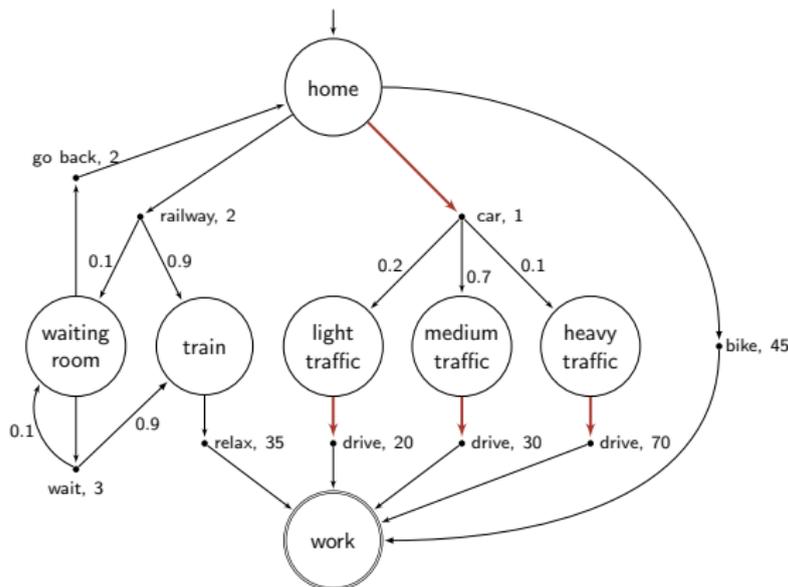
- ▷ What kind of **strategies** are we looking for when the environment is **stochastic** (MDP)?

## Solution 1: minimize the *expected* time to work



- ▷ “Average” performance: meaningful when you journey often.
- ▷ **Simple strategies** suffice: no memory, no randomness.
- ▷ Taking the **car** is optimal:  $\mathbb{E}_D^\sigma(TS^{\text{work}}) = 33$ .

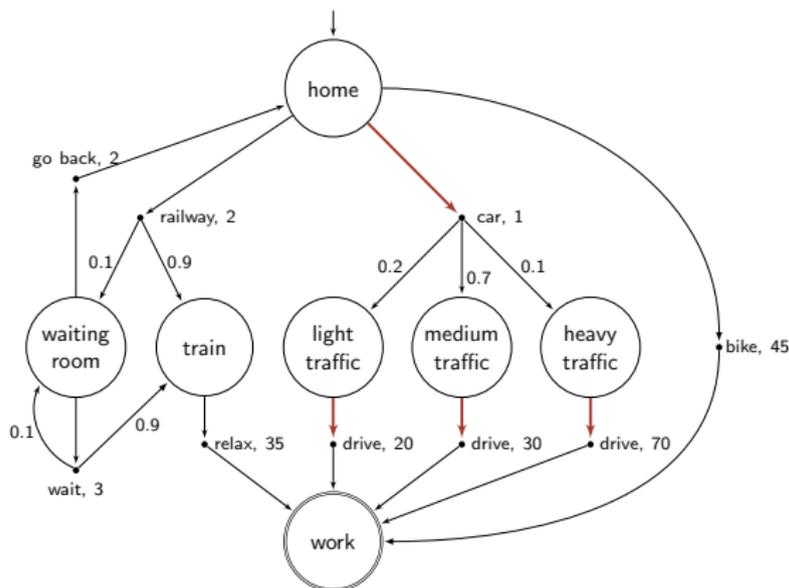
## Solution 2: traveling without taking too many risks



Minimizing the *expected time* to destination makes sense **if** we travel often and **it is not a problem to be late**.

With car, in 10% of the cases, the journey takes 71 minutes.

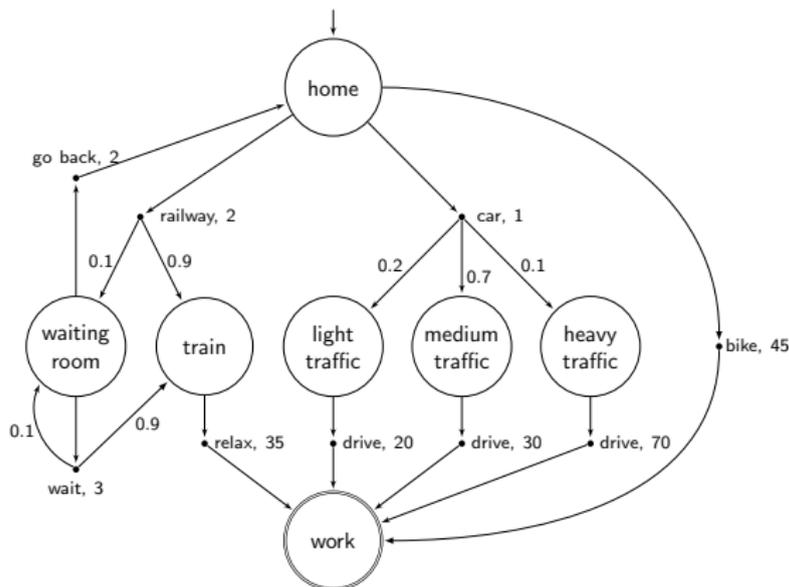
## Solution 2: traveling without taking too many risks



**Most bosses will not be happy if we are late too often...**

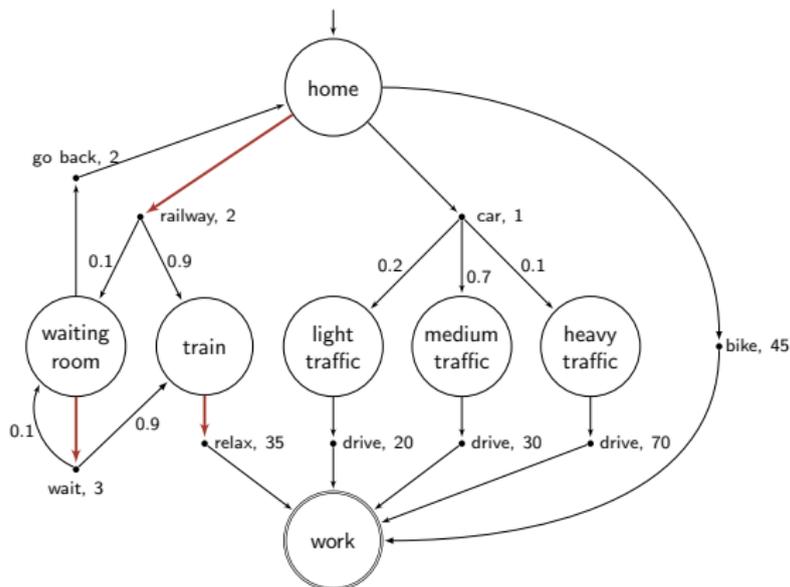
~> what if we are risk-averse and want to avoid that?

## Solution 2: maximize the *probability* to be on time



**Specification:** reach work within 40 minutes with 0.95 probability

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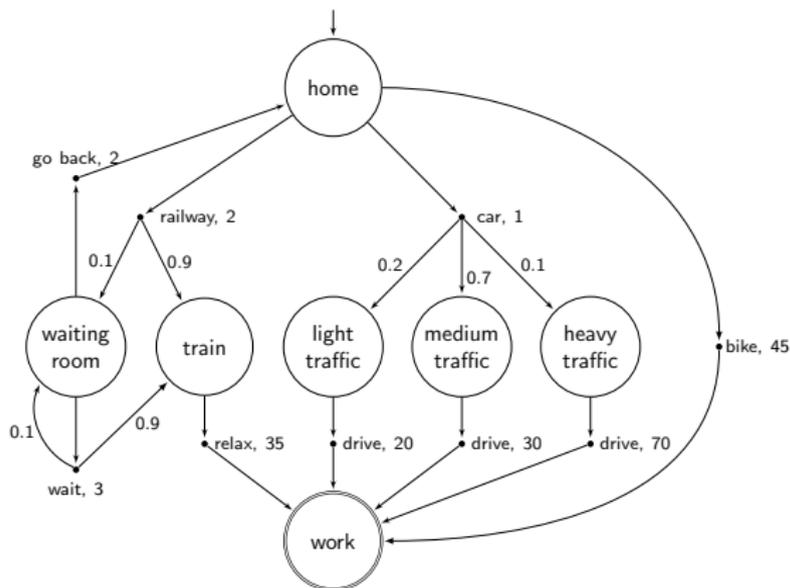


**Specification:** reach work within 40 minutes with 0.95 probability

**Sample strategy:** take the **train**  $\rightsquigarrow \mathbb{P}_D^\sigma [\text{TS}^{\text{work}} \leq 40] = 0.99$

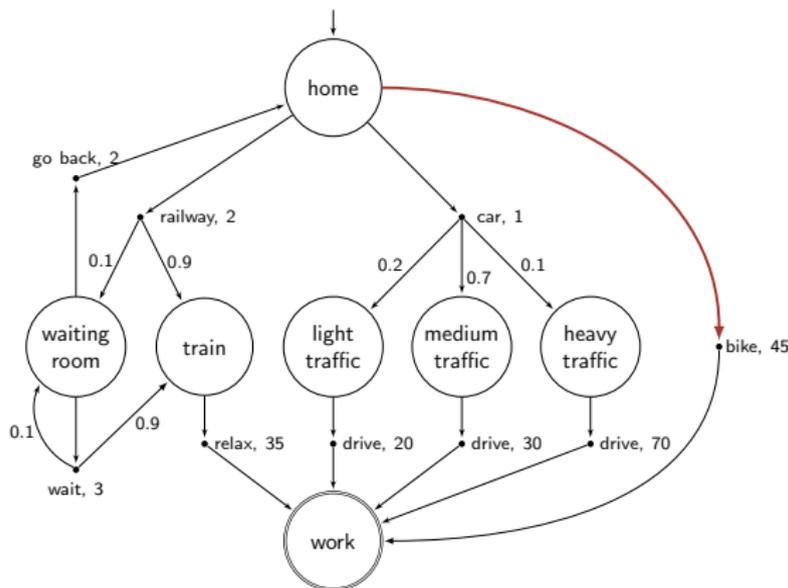
**Bad choices:** car (0.9) and bike (0.0)

## Solution 3: strict worst-case guarantees



**Specification:** *guarantee* that work is reached within 60 minutes (to avoid missing an important meeting)

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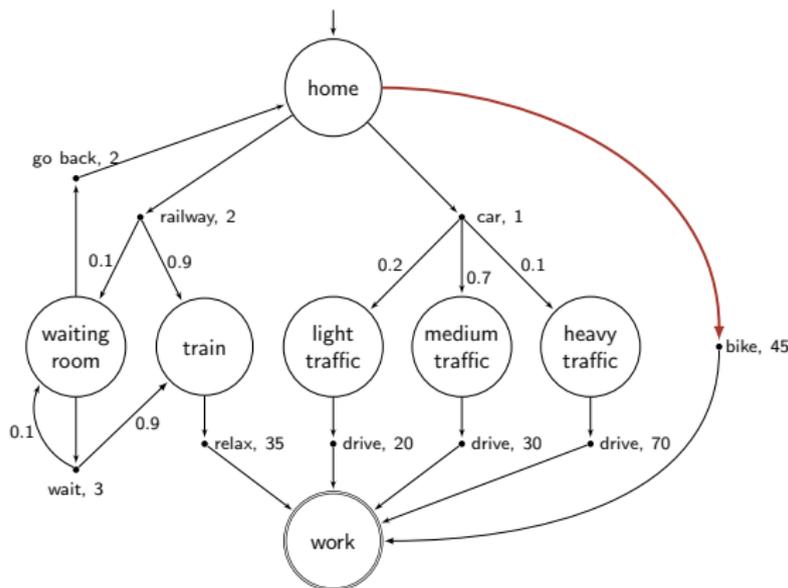


**Specification:** *guarantee* that work is reached within 60 minutes (to avoid missing an important meeting)

**Sample strategy:** **bike**  $\rightsquigarrow$  worst-case reaching time = 45 minutes.

**Bad choices:** train ( $wc = \infty$ ) and car ( $wc = 71$ )

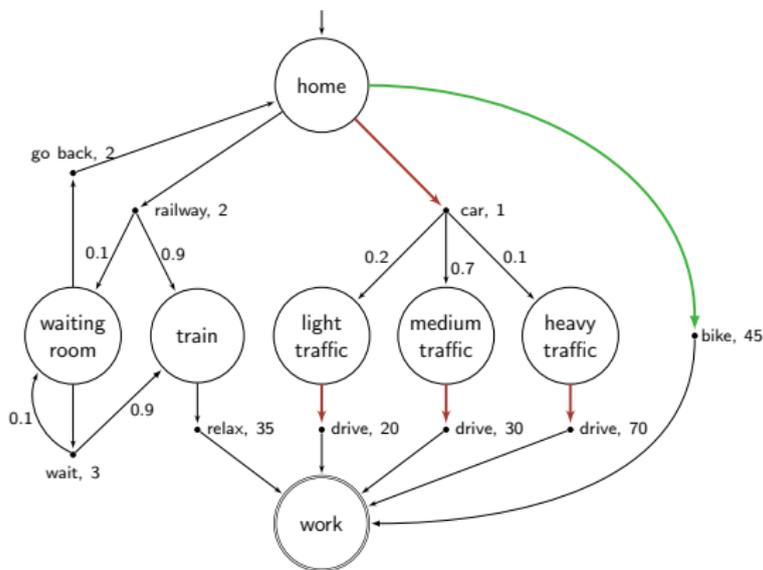
## Solution 3: strict worst-case guarantees



Worst-case analysis  $\rightsquigarrow$  **two-player game** against an antagonistic adversary (*bad guy*)

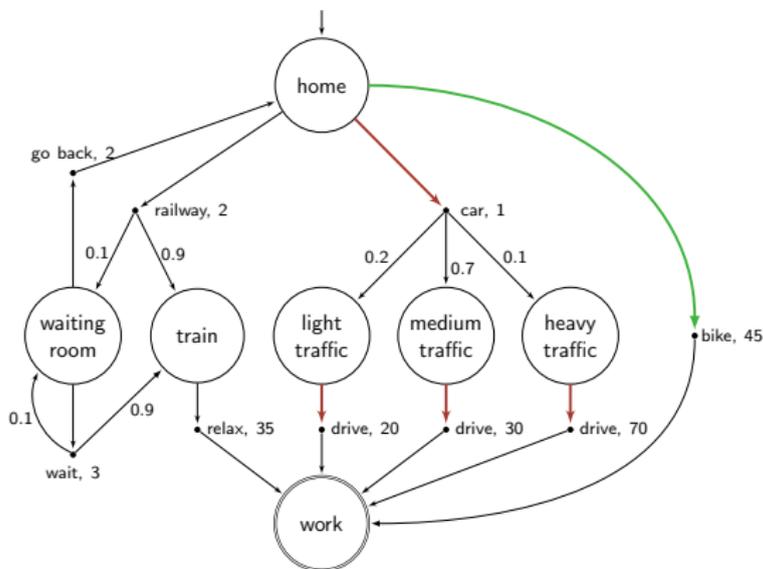
- ▷ forget about probabilities and give the choice of transitions to the adversary

## Solution 4: minimize the *expected* time under strict worst-case guarantees



- Expected time: **car**  $\rightsquigarrow \mathbb{E} = 33$  but **wc** = 71 > 60
- Worst-case: **bike**  $\rightsquigarrow$  wc = 45 < 60 but  $\mathbb{E} = 45 \gg \gg 33$

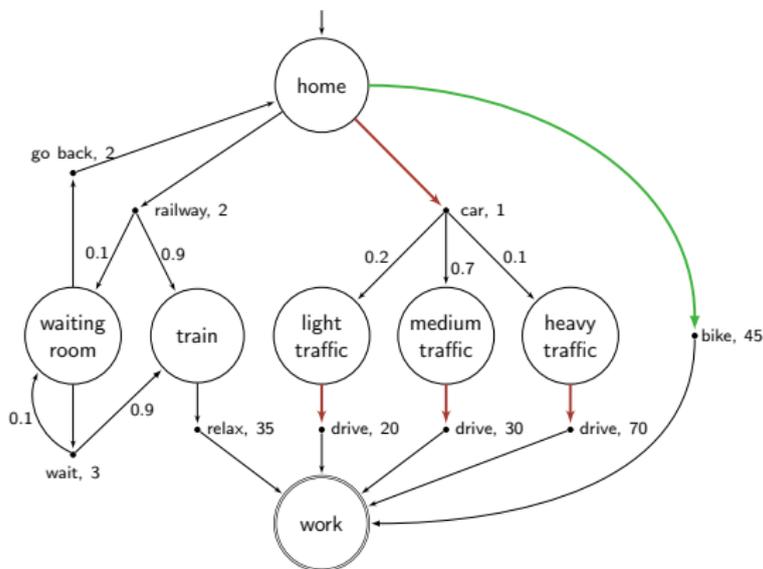
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In practice, we want both! Can we do better?

- ▶ **Beyond worst-case synthesis** [BFRR14b, BFRR14a]: minimize the expected time under the worst-case constraint.

## Solution 4: minimize the *expected* time under strict worst-case guarantees

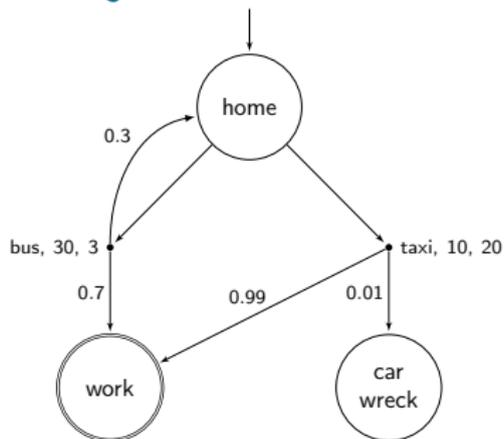


**Sample strategy:** try train up to 3 delays then switch to bike.

↪  $wc = 58 < 60$  and  $\mathbb{E} \approx 37.34 \ll 45$

↪ Strategies need **memory** ↪ more complex!

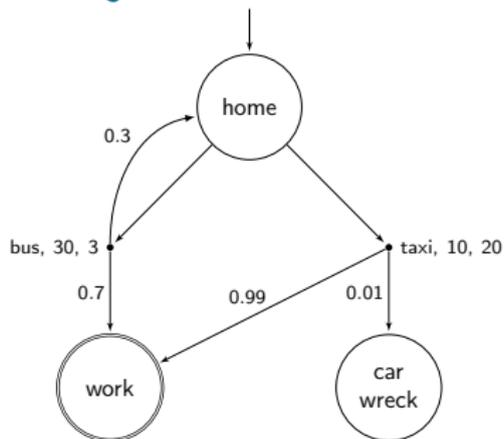
## Solution 5: multiple objectives $\Rightarrow$ trade-offs



Two-dimensional weights on actions: *time* and *cost*.

Often necessary to consider **trade-offs**: e.g., between the probability to reach work in due time and the risks of an expensive journey.

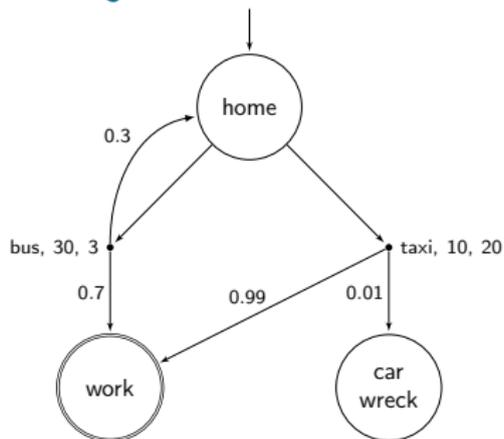
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Solution 2 (probability) can only ensure a **single constraint**.

- **C1:** 80% of runs reach work in at most 40 minutes.
  - ▷ Taxi  $\rightsquigarrow \leq 10$  minutes with probability  $0.99 > 0.8$ .

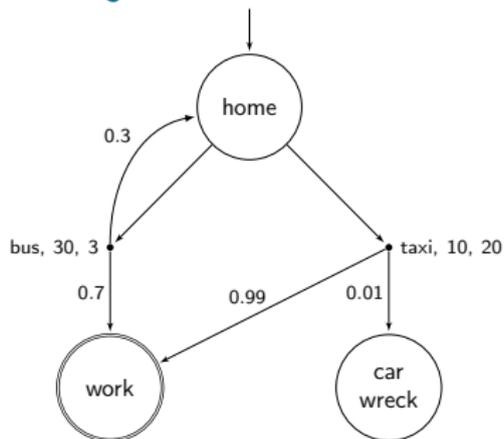
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  - ▷ Bus  $\rightsquigarrow \geq 70\%$  of the runs reach work for 3\$.

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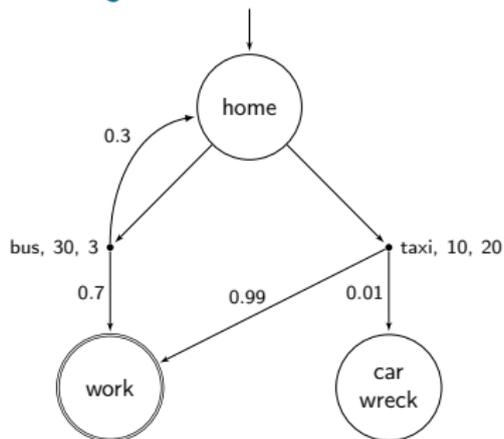


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Taxi  $\not\models$  C2, bus  $\not\models$  C1. What if we want  $C1 \wedge C2$ ?

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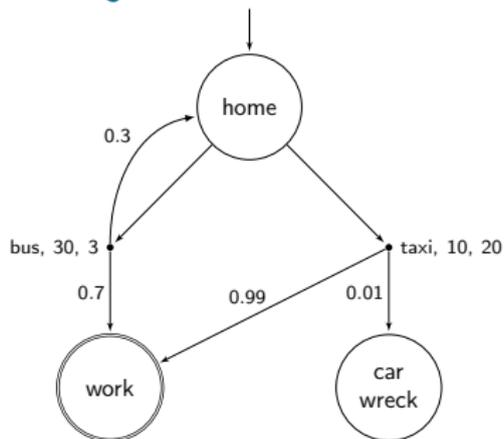


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Study of **multi-constraint percentile queries** [RRS15a].

- ▷ Sample strategy: bus once, then taxi. Requires *memory*.
- ▷ Another strategy: bus with probability 3/5, taxi with probability 2/5. Requires *randomness*.

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- **C1:** 80% of runs reach work in at most 40 minutes.
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In general, *both memory and randomness* are required.

$\neq$  previous problems  $\leadsto$  more complex!

## Conclusion (1/2)

This talk was about **shortest path objectives**, but there are many more! Some examples based on energy applications.

- ▶ **Energy**: operate with a (bounded) fuel tank and never run out of fuel [BFL<sup>+</sup>08].
- ▶ **Mean-payoff**: average cost/reward (or energy consumption) per action in the long run [EM79].
- ▶ **Average-energy**: energy objective + optimize the long-run average amount of fuel in the tank [BMR<sup>+</sup>15].

## Conclusion (2/2)

Our research aims at:

- defining meaningful *strategy concepts*,
- providing *algorithms* and *tools* to compute those strategies,
- classifying the *complexity* of the different problems from a theoretical standpoint.
  - ↪ Is it mathematically possible to obtain efficient algorithms?

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Thank you! Any question?

# References I



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